



Capacitated Vehicle Routing Problem Solving through Adaptive Sweep based Clustering plus Swarm Intelligence based Route Optimization

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Abstract

Capacitated Vehicle Routing Problem (CVRP) is an optimization task where customers are assigned to vehicles aiming that combined travel distances of all the vehicles as minimum as possible while serving customers. A popular way among various methods of CVRP is solving it in two phases: grouping or clustering customers into feasible routes of individual vehicles and then finding their optimal routes. Sweep is well studied clustering algorithm for grouping customers and different traveling salesman problem (TSP) solving methods are commonly used to generate optimal routes of individual vehicles. This study investigates effective CVRP solving method based on recently developed adaptive Sweep and prominent Swarm Intelligence (SI) based TSP optimization methods. The adaptive Sweep cluster is a heuristic based adaptive method to select appropriate cluster formation starting angle of Sweep. Three prominent SI based TSP optimization methods are investigated which are Ant Colony Optimization, Producer-Scrounger Method and Velocity Tentative Particle Swarm Optimization (VTPSO). Genetic Algorithm is also considered since it is the pioneer and well-known population based method. The experimental results on two suites of benchmark CVRPs identified the effectiveness of adaptive Sweep plus SI methods in solving CVRP. Finally, adaptive Sweep plus the VTPSO is found better than other tested methods in this study as well as several other prominent existing methods.



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Introduction

Vehicle Routing Problem (VRP) is a complex combinatorial optimization task and has been widely studied since introduced by Dantzig and Ramser in

1959¹⁻⁶. VRP can be described as the problem of designing optimal delivery or collecting routes from one or several depots to a number of geographically scattered customers, subject to different constraints.

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Capacitated VRP (CVRP) is the most general form of VRP with an additional constraint of fixed vehicle capacity⁷⁻¹¹.

CVRP is a one of the most studied problems which works with predefined demands and locations of customers to serve with fixed number of vehicles. It constructs routes of the vehicles in such a way that: (i) every route starts and ends at the depot; (ii) all the demands are accomplished; (iii) the vehicle's capacity is not exceeded; (iv) a customer is visited by only a single vehicle; (v) the sum of costs is low as possible. The aim of CVRP solving is to minimize the combined traveling distance or time for all vehicles while serving all the customers.

Various CVRP solving methods have been proposed recently. A number of methods optimizes vehicles' customer assignment and vehicles' route generation together¹². Otherwise, grouping or clustering of customers into feasible routes maintaining given constraints and then finding optimal routes of vehicles is the most effective way of solving CVRP⁴. Sweep algorithm is the most popular clustering method among several ways of grouping customers and is well-studied due to its simplicity. The method creates clusters based on the angular position of the customers calculating polar angles of all their positions^{12,13}. Inserting customers into a cluster is done in either clock-wise (i.e., backward Sweep) or anti clock-wise (i.e., forward Sweep) direction until all the customers are visited¹⁴. On the other hand, Traveling Salesman Problem (TSP) optimization methods are generally used to find optimal route of each individual vehicle^{14,15}.

Some studies are available to solve specific CVRP tasks using Sweep clustering and TSP optimization methods. Nurcahyo *et al.*,¹⁴ investigated public transport of Semarang, Indonesia using Sweep based VRP. They considered nearest neighbor algorithm of TSP for route generation. Suthikarnnarunai¹⁶ solved routing problem of a University of Bangkok using Sweep algorithm for clustering and TSP routes were generated through integer programming. Aziz *et al.*,¹⁰ investigated a hybrid algorithm of Sweep and nearest neighbor algorithm to solve CVRP. They tested the method on Augerat's Euclidean benchmark dataset and in solving the dairy products delivery problem of Tiba Company for Trade and Distribution in Egypt.

Na *et al.*,¹⁷ introduced nearest neighbor in Sweep and optimized route using 2-opt edge exchange method. Author made several extensions of sweep algorithm but an extra operation increases the computation cost.

Genetic Algorithm (GA) and different Swarm Intelligence (SI) based methods are applied in route optimization on cluster formation using Sweep algorithm. Nazif¹⁸ investigated optimized crossover GA for solving CVRP. Yousefikhoshbakht¹⁹ proposed a hybrid algorithm combining Ant Colony System (ACO), Sweep algorithm and 3-opt local search for solving CVRP. Reed *et al.*,¹² demonstrated gathering of reusing waste from family units using ACO. Tan *et al.*,²⁰ used several heuristics combined with ACO to solve CVRP and showed it a viable alternative to solve CVRP. Kao *et al.*,²¹ proposed a new hybrid algorithm based on ACO and Particle Swarm Optimization (PSO) for solving CVRP. Kanthavel and Prasad²² investigated Nested PSO (NPSO) for route generation. Tavakoli and Sami²³ and Venkatesan *et al.*,²⁴ are also considered in CVRP. Pornsing²⁵ proposed two novel PSO-based algorithms to solve CVRP named Survival Sub-swarms Adaptive PSO (SSS-APSO) and Survival Sub-swarms Adaptive PSO with velocity-line bouncing (SSS-APSO-vb).

The objective of this study is to investigate effective CVRP solving technique through recently developed adaptive Sweep²⁷ and prominent SI based TSP optimization methods. The adaptive Sweep cluster is a heuristic based adaptive method to select appropriate cluster formation starting angle of Sweep. Three prominent SI based TSP route optimization methods are used in this paper which are Ant Colony Optimization (ACO), Producer-Scrounger Method (PSM) and Velocity Tentative PSO (VTPSO). GA is also considered since it is the pioneer and well-known population based method. The adaptive Sweep is the expansion of standard Sweep selecting cluster formation starting angle adaptively²⁷. In conventional Sweep algorithm, grouping customers begins from 0° and thusly progresses toward 360° to relegate every one of the customers under various vehicles¹⁶. Because of such unbending beginning from 0° total number of clusters may surpasses total number of accessible vehicles for a few examples. Beginning from various angles subsequently accomplished better CVRP result²⁶. The outcome of adaptive

Sweep plus SI based methods are experimented on a huge number of benchmark CVRPs and identified the effectiveness comparing outcomes with other existing techniques.

The overview of the paper is as per the followings. Section 2 describes the method of solving CVRP using adaptive Sweep and SI methods (GA, ACO, PSM and VTPSO) briefly. Section 3 reports the experimental results on the benchmark problems considering adaptive Sweep clustering with each of GA, ACO, PSM and VTPSO. Finally, Section 4 concludes the paper with few remarks.

CVRP Solving through Adaptive Sweep based Clustering plus SI based Route Optimization

This section first gives description of adaptive Sweep with deficiency of standard Sweep and then describes considered TSP methods to make the paper self-contained as well as better understanding.

Adaptive Sweep Clustering

Gillett and Miller¹³ coined the name “the sweep algorithm” as a heuristic algorithm. Initial routes or clusters are formed by sweeping the nodes according to their polar angle (increasing or decreasing order). Sweeping halts when vehicle capacity constraint is violated, finishes a single vehicle route and resumes for another vehicle. Cluster formation starts in standard Sweep with an arbitrary customer (from 0°) and then sequentially assigns the remaining customers (moving toward 360°) to assign all the customers under available vehicles^{12,15,16,28}. In several cases such type of clustering produces total number of clusters more than total number of vehicles. To overcome the problem, adaptive Sweep²⁷ heuristically identifies the appropriate starting angle (Θ_s) of cluster formation for any given instance. The method first computes polar angle of customers and order those to a list (say ONL) according to polar angle. The approach considers angle difference of consecutive nodes in ONL; and distance between the nodes and distances from the depot. Then preference value ($p\Theta$) of each consecutive nodes is calculated and cluster formation starts from maximum $p\Theta$ value. For example the depot and other two consecutive customers are D, C1 and C2, respectively. Polar angles of the customers are Θ_1 and Θ_2 . The distances of the customers from the depot is dC_1 and dC_2 ; and distance between

the customers is dC_{12} . Preference value ($p\Theta$) for the starting angle between the customers C1 and C2 means to place the customers in two different clusters and is calculated using Eq. (1).

$$p\Theta = \alpha * (\Theta_2 - \Theta_1) + \beta * \{dC_{12} + \text{Min}(dC_1, dC_2)\} \quad \dots(1)$$

In the equation, α and β are the arbitrary constants to emphasis angle difference and node distances, respectively.

Algorithm 1 shows the steps of adaptive Sweep algorithm. First three steps of the initialization section are same as standard Sweep: update nodes' coordinates considering depot location as (0,0), compute polar angle of each customer and order the nodes according to polar angle to a list ONL. The main significance of adaptive Sweep is that it starts cluster formation from the maximum preference values. First the method calculates starting angle of cluster formation ($p\Theta$) according to Eq. (1) (Step 1.e). As like standard Sweep, the method assigns nodes into a cluster while vehicle capacity does not exceed (Steps 2.b and 2.d) otherwise new cluster forms for unassigned nodes (Step 2.e). Since the adaptive Sweep may starts any location of ONL, Step 2.e transforms node assignment from bottom of ONL to the beginning of ONL. It is notable that for $\Theta_s = 0^\circ$ the proposed method will be standard Sweep.

Algorithm 1: Adaptive Sweep Algorithm

Initialization

- Calculate co-ordinates of the customers considering depot as (0, 0).
- Calculate the polar angle of each customer.
- Order the customers according to polar angle, ONL
- Calculate distance between the customers and distances from the depot.
- Calculate preference value ($p\Theta$) of each consecutive nodes using Eq. 1.

Clustering

- Cluster $C=1$.
- Take maximum preference value as starting angle of cluster formation , $p\Theta$
- Add customers to current cluster C.
Stop when including the next customer would

- exceed vehicle capacity.
- Make another cluster C+1 by continuing the sweep where the previous customer left off.
- Repeat Steps 2.b - 2.d, until all customers have been incorporated in a cluster.

Outcome

All the customers are relegated into total C clusters

Optimal Route Generation of Vehicles

Optimal route generation of each individual vehicle is a crucial part of CVRP solving while any clustering method is used to cluster customers. In general, a clustering method divides total CVRP nodes into clusters²⁹, whereby number of clusters is equal to the number of vehicles. The aim of route generation is the optimal path finding of each vehicle starting from the depot and returning to depot after serving all of its assigned nodes. Therefore, route generation of individual-vehicle is simply a small sized TSP considering the depot as a common city point; and any TSP optimization method may be used for this purpose. To generate route for a vehicle, a TSP cost matrix considering nodes for a particular vehicle is prepared and then a TSP optimization is employed to work with the cost matrix as an independent TSP. Following sub-sections briefly describes TSP methods considered in this study which are GA, ACO, PSM and VTPSO.

Genetic Algorithm (GA)

GA is one of the most popular search and optimization techniques based on the natural evolution through genetic inheritance. It works with populations of chromosomes, selection according to fitness, crossover to produce new offspring, and random mutation of new offspring.

Selection operation selects good solutions in a population and forms a mating pool. A number of selection techniques are used in GAs. Roulette Wheel Selection, Random Selection, Rank Selection and Tournament Selection are mainly used to select the parents in GA.

A crossover operator is used in GA to recombine two solutions to get a better solution. Crossover in biological terms refers to the blending of chromosomes from the parents to produce new

chromosomes for the offspring. Two strings are picked from the mating pool at random to crossover³⁰. Among several crossover techniques, Enhanced Edge Recombination (EER) method is used to solve TSP. In EER, an adjacency table³⁰ (called Edge Table) is prepared that lists links into and out of a city found in the two parent sequences. Element of a sequence with a common edge is marked as inverting sign to emphasis in selection. The description of EER is available in³⁰.

Mutation³⁰ is the process by which offsprings are generated with a single parent. Position swap of two randomly selected nodes is the common way of mutation operation for TSP.

Elitism saves the best chromosome to the new offspring population before crossover and mutation to eliminate lose of best chromosome. Elitism keeps the best solutions to a stack and helps to improve performance of GA.

Ant Colony Optimization (ACO)

ACO was inspired by behaviors of real ants³¹ while searching food. Initially, ants choose different paths if there exists several paths between ant colony and the food source. After sometimes all of them follow the shortest path; it is because of pheromone. Pheromone is a chemical that ants lay in their path. More pheromone means more ants travelled the path and also means the path is comparatively shorter. The general ACO is relatively simple and based on a set of ants, each making one of the possible round-trips along the cities. If an ant in city i , the probability to go city j can be calculated by the following equation and parameters:

$$P_{i,j}^k(t) = \frac{[\tau_{i,j}(t)]^\alpha [\eta_{i,j}]^\beta}{\sum_{i \in J_i^k} [\tau_{i,j}(t)]^\alpha [\eta_{i,j}]^\beta}, \quad \dots(2)$$

J_i^k is the set of cities the ant still has to visit.

$\eta_{i,j} = 1/d_{i,j}$ is the reciprocal of the distance from i to j .

$\tau_{i,j}$ is the amount of pheromone on the arc from i to j .

α is the importance of the intensity in the probabilistic transition.

β is the importance of the visibility of the trail segment.

After the completion of a tour, each ant lays some pheromone on the path. The pheromone is updated by the following equations.

$$\forall (i, j) \tau_{ij}(t+1) = (1 - \rho) \cdot \tau_{ij} + \sum_{k=1}^m \Delta \tau_{ij}^k(t) \quad \dots(3)$$

$$\Delta \tau_{ij}^k = \begin{cases} \frac{1}{L_k} & \text{if ant } k \text{ uses edge } i, j \\ 0 & \text{otherwise} \end{cases} \quad \dots(4)$$

ρ is the trail persistence or evaporation rate. The detail description of ACO available in³¹.

Producer Scrounger Method (PSM)

PSM³² is a TSP solving method which is inspired from the collective behavior of animal group. It models roles and cooperation of three classes of animal group members: producer, scrounger and dispersed. Here producers have the best tour, few dispersed members have worse tours and they randomly checks new tours. At each step of PSM, the producer searches better tour, scroungers traverse new tours while moving toward producer's tour; and dispersed members arbitrarily examines new tours. In case of making producer's tour, PSM arbitrarily selects a city from the producer's tour and exchanges its connection with other closest cities for better tours. Parameter rate of near cities (RNC) defines the number of cities to be checked by the producer for better tour. A scrounger updates position towards the producer through Swap operator and swap sequence. A Swap Operator demonstrates two cities in a tour those positions will be exchanged. Suppose, a TSP problem has five cities and a solution is A-B-C-D-E. A Swap Operator SO(1,2) gives the new solution S'.

$$\begin{aligned} S' &= S + \text{SO}(1,2) \\ &= (\text{A-B-C-D-E}) + \text{SO}(1,2) \\ &= \text{B-A-C-D-E} \end{aligned} \quad \dots(5)$$

A swap sequence is formed from one or more swap operators. Finally, producer is considered as the solution of a given problem. The method performs well when tested on a suite of benchmark TSPs. The detailed description of this method is available in³².

Velocity Tentative PSO (VTPSO)

VTPSO³³ is the most recent SI based method extending PSO to solve TSP. It considers Swap Sequence (SS) as velocity and calculates as like

conventional standard SS based PSO method³⁴ but apply the SS in a different way. In traditional PSO, the new tour is calculated implementing all the SOs of a SS without considering no intermediate tours. But VTPSO conceives a measure (called partial search) to apply such velocity to change particles position. VTPSO measures tours with portions of SS and conceive comparatively better new tour with a portion or full tentative SS. It calculates velocity SS using Eq. (6) considering (i) last applied velocity ($v^{(t-1)}$), (ii) previous best solution of the particle (P) and (iii) global best solution of the swarm (G).

$$V_i^{(t)} = V_i^{(t-1)} \otimes \alpha (P_i - X_i^{(t-1)}) \otimes \beta (G - X_i^{(t-1)}) \quad \alpha, \beta \in [1,0] \quad \dots(6)$$

The tentative tour $X_{i(t)}$ is calculated using Eq. (7) having the minimum tour cost.

$$X_i^{(t)} = X_i^{j(t)} \quad \dots(7)$$

The method is shown to perform well when tested on a suite of benchmark TSPs. The detailed description of this method is available in³⁴.

Experimental Results

This section checks adequacy of Adaptive Sweep algorithm and SI methods in solving benchmark CVRPs. Description of the benchmark problems and experimental setting are explained first.

Bench Mark Datasets and Experimental Settings

Two different sets of Augerat benchmark problems³⁴ (A-VRP and P-VRP) have been considered in this study. In A-VRP, number of customers varies from 32 to 80, total demand varies from 407 to 942, number of vehicles varies from 5 to 10 and capacity of individual vehicle is 100 for all the problems. For example, A: n32-k5 has 32 customers and 5 vehicles. On the other hand, in P-VRP, number of customers varies from 16 to 101, total demand varies from 246 to 22500 and vehicle capacity varies from 35 to 3000. Table 1 and Table 2 depict the brief description of the A-VRP and P-VRP benchmark problems, respectively. The numeric value in a problem name presents the number of customer nodes and vehicles. The detailed description of the problems are available in provider's website³⁵. According to Table 1 and Table 2, the selected benchmark problems belongs large verities in number of nodes,

vehicles and demands; and therefore, provides a diverse test bed.

A customer node is represented as a co-ordinate in a problem. Standard Sweep starts cluster formation from 0° (i.e., $\Theta_s = 0^\circ$) and does not have any parameter to set. In adaptive Sweep, the values of α and β were set to 0.6 and 0.2, respectively and found effective for most of the problems; and tuned

between 0.2 and 0.6 for few other problems. In ACO, alpha and beta were set to 1 and 3, respectively. On the other hand, the RNC (rate of near cities consideration) for producer scanning in PSM was set to 0.1. The algorithms are implemented on Visual C++ of Visual Studio 2013. The experiments have been done on a PC (Intel Core i5-3470 CPU @ 3.20 GHz CPU, 4GB RAM) with Windows 7 OS.

Table 1: Augerat A-VRPs CVRP problems

SI.	Name of Problem	No. of Nodes	Number of Vehicles	Vehicle Capacity	Demand
1	n32-k5	32	5	100	410
2	n33-k5	33	5	100	446
3	n33-k6	33	6	100	541
4	n34-k5	34	5	100	460
5	n36-k5	36	5	100	442
6	n37-k5	37	5	100	407
7	n37-k6	37	6	100	570
8	n38-k5	38	5	100	481
9	n39-k5	39	5	100	475
10	n39-k6	39	6	100	526
11	n44-k6	44	6	100	570
12	n45-k6	45	6	100	593
13	n45-k7	45	7	100	634
14	n46-k7	46	7	100	603
15	n48-k7	48	7	100	626
16	n53-k7	53	7	100	664
17	n54-k7	54	7	100	669
18	n55-k9	55	9	100	839
19	n60-k9	60	9	100	829
20	n61-k9	61	9	100	885
21	n62-k8	62	8	100	733
22	n63-k9	63	9	100	873
23	n63-k10	63	10	100	932
24	n64-k9	64	9	100	848
25	n65-k9	65	9	100	877
26	n69-k9	69	9	100	845
27	n80-k10	80	10	100	942

Table 2: Augerat P-VRP's CVRP problems

SI.	Problem Name	No. of Nodes	Number of Vehicles	Vehicle Capacity	Demand
1	n16-k8	16	8	35	246
2	n19-k2	19	2	160	310
3	n20-k2	20	2	160	310

4	n21-k2	21	2	160	298
5	n22-k2	22	2	160	308
6	n22-k8	22	8	3000	22500
7	n23-k8	23	8	40	313
8	n40-k5	40	5	140	618
9	n45-k5	45	5	150	692
10	n50-k7	50	7	150	951
11	n50-k8	50	8	120	951
12	n50-k10	50	10	100	951
13	n51-k10	51	10	80	777
14	n55-k7	55	7	170	1042
15	n55-k8	55	8	160	1042
16	n55-k10	55	10	115	1042
17	n55-k15	55	15	70	1042
18	n60-k10	60	10	120	1134
19	n60-k15	60	15	80	1134
20	n65-k10	65	10	130	1219
21	n70-k10	70	10	135	1313
22	n76-k4	76	4	350	1364
23	n76-k5	76	5	280	1364
24	n101-k4	101	4	400	1458

Detailed Experimental Observation on Selected Problems

This section shows the detailed results for the selected benchmark problems A: n53-k7 and P: n65-k10 of A-VRP and P-VRP. For GA, PSM and VTPSO population size was 100; whereas, number of ants in ACO equals the number of customers assigned to a vehicle as it desire. The iteration number is set at 200 for the algorithms. Table 3 shows the total clusters for different fixed as well as adaptively selected starting angles (Θ_s) and optimized route cost with different methods for A: n53-k7 problem. The problem has 53 nodes and total 664 demand to be served with seven vehicles having capacity 100. From the table it is observed that total number of clusters for

$\Theta_s=0^\circ$ (i.e., in standard Sweep) is 8 that is more than available vehicles. Total clusters are also 8 for $\Theta_s=270^\circ$. On the other hand, number of clusters equal to total vehicles (i.e., 7) for $\Theta_s=90^\circ$ and 180° . It is also remarkable that CVRP cost (i.e., total travel distance) for 7 clusters is lower than the cases of 8 clusters after route optimization. It is interesting from the table that total clusters are also 7 for adaptively selected angle 220.6° . The best CVRP cost for an algorithm among different Θ_s is marked as bold-faced type. For the problem the best CVRP cost achieved after optimizing for with GA, ACO, PSM and VTPSO are 1091, 1131, 1190 and 1090, respectively. The best values are found for adaptively selected $\Theta_s = 220.6^\circ$ and fixed $\Theta_s=180^\circ$.

Table 3: Outcome comparison using GA, ACO, PSM and VTPSO for A: n53-k7 problem of A-VRP

Θ_s	Clusters	CVRP Cost Before Route Optimizing	CVRP cost after optimizing with			
			GA	ACO	PSM	VTPSO
0°	8	1604	1174	1211	1174	1174
90°	7	1654	1125	1160	1109	1109
180°	7	1504	1091	1131	1090	1090
270°	8	1775	1171	1196	1171	1171
220.6^{0*}	7	1504	1091	1131	1090	1090

* Starting angle selected through adaptive Sweep

Table 4 shows the total clusters for different fixed as well as adaptively selected starting angles (Θ_s) and optimized route cost with different methods for P: n65-k10 problem of P-VRP. The problem has 65 nodes and total 1219 demand to be served with ten vehicle having capacity 130. From the table it is observed that total number of clusters for $\Theta_s=0^\circ$ (i.e., in standard Sweep) is 11 that is more than available vehicles. Total clusters are also 11 for $\Theta_s=90^\circ$. On the other hand, number of clusters equal to total

vehicles (i.e., 10) for $\Theta_s=180^\circ$ and 270° . It is also remarkable that final CVRP cost for 10 clusters is lower than the cases of 11. For the problem the best CVRP cost achieved (i.e., 837) for $\Theta_s=180^\circ$ with GA, PSM and VTPSO. On the other hand, the heuristic approach selected starting angle is $\Theta_s = 278.43^\circ$ and outcome is same for fixed $\Theta_s = 270^\circ$ with 10 clusters. Although the outcome is inferior to best outcome with $\Theta_s = 180^\circ$, the outcome is better than standard Sweep with $\Theta_s = 0^\circ$.

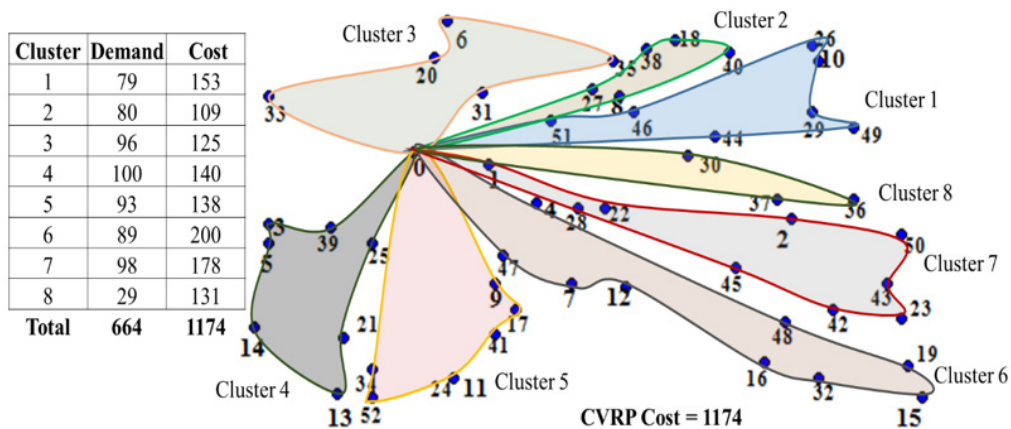
Table 4: Outcome comparison using GA, ACO, PSM and VTPSO for P: n65-k10 problem of P-VRP

Θ_s	Clusters	CVRP Cost Before Route Optimizing	CVRP cost after optimizing with			
			GA	ACO	PSM	VTPSO
0°	11	1142	864	933	864	864
90°	11	1151	877	946	874	874
180°	10	1154	837	900	837	837
270°	10	1256	860	890	859	859
278.43^{0*}	10	1256	860	890	859	859

* Starting angle selected through adaptive Sweep

The graphical representation of the solution of A: n53-k7 problem for standard Sweep clustering (i.e., $\Theta_s=0^\circ$) is shown in Fig. 1. Eight clusters are generated and Cluster 8 is for unassigned three nodes having total demand 29. Otherwise, Cluster 1 covers total demand of 79 although vehicle capacity 100. The CVRP costs for route optimization with GA and ACO are 1174 and 1212, respectively. On the

other hand, PSM and VTPSO gave same solution with CVRP cost 1174 as shown in Fig. 1(a). The reason for worst CVRP cost with ACO might be inclination with pheromone in ACO and solutions for Cluster 4 and Cluster 6 are bad with respect to other methods. On the other hand, slightly different solution of GA from PSM/VTPSO is shown for Cluster 6.



(a) Route optimization using GA, PSM or VTPSO

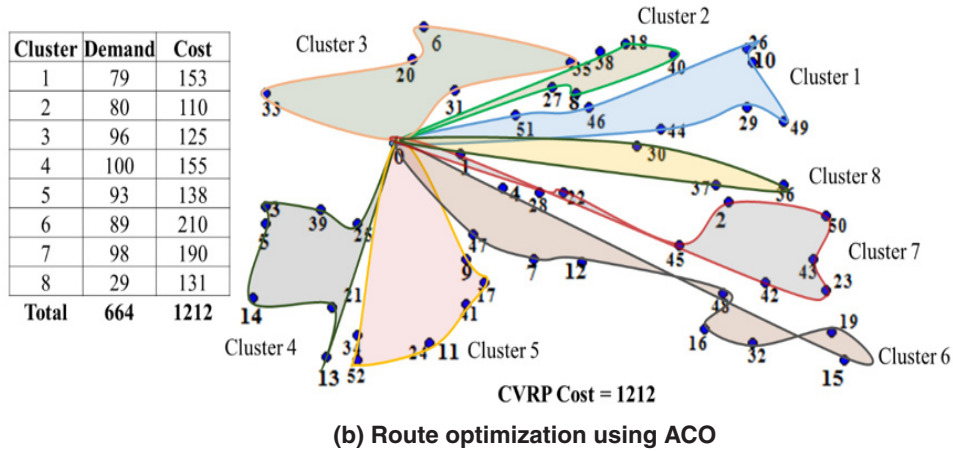
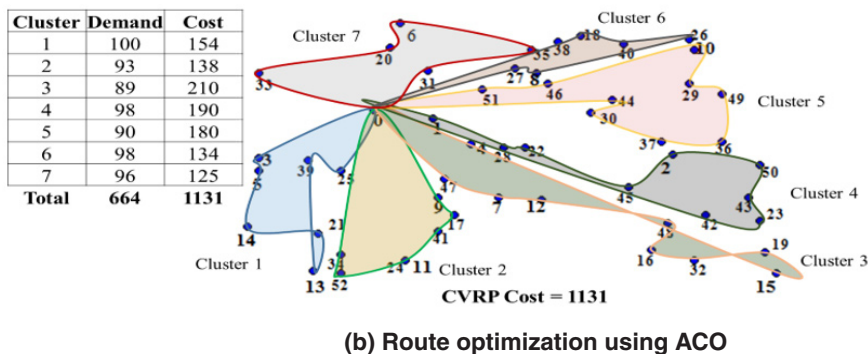
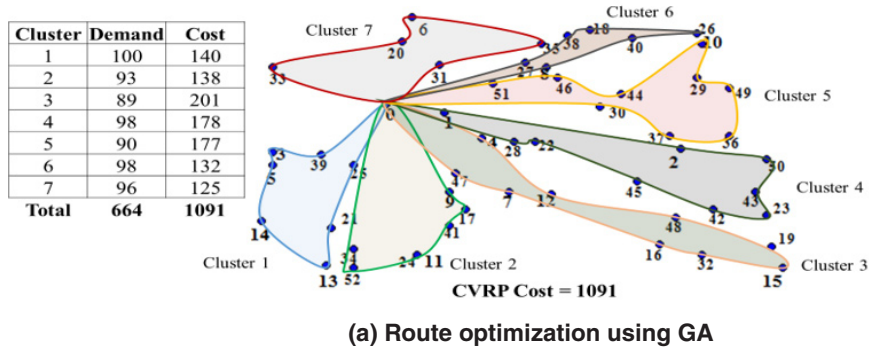
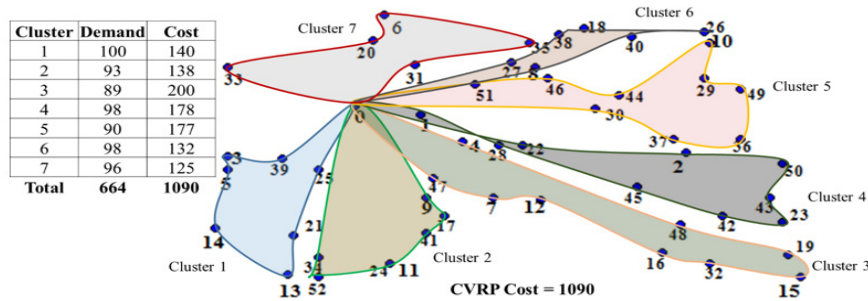


Fig. 1: Pictorial view of solutions with standard Sweep clustering (i.e., $\Theta_s=0^0$) for A: n53-k7 problem

Figure 2 shows the graphical representation of adaptive sweep clustering for A: n53-k7 problem with adaptively selected $\Theta_s = 220.6^0$. In this case total demands are fulfilled by seven clusters that is equal to number of vehicles. Among the four route optimization methods, CVRP cost with ACO is the worst and the value is 1131. Similar to standard Sweep, it achieved worse solution for Cluster 4 and Cluster 6. The best CVRP solution for the

problem is achieved by PSM and VTPSO and achieved CVRP cost is 1090. On the other hand, GA is showed competitive result with PSM/VTPSO showing different result only for Cluster 6 and CVRP cost 1091. Finally, the comparative description with graphical representation of Figs. 1 and 2 clearly identified the proficiency of adaptive Sweep over standard Sweep.





(c) Route optimization using PSM and VTPSO

Fig. 2: Pictorial view of solutions with adaptive Sweep clustering with adaptively selected $\Theta_s = 220.60$ for A: n53-k7 problem

CVRP Outcomes and Performance Comparison

This section shows the proficiency of adaptive Sweep clustering over standard Sweep clustering while using GA, ACO, PSM and VTPSO for route optimization. Finally the outcome of the adaptive sweep with the prominent methods in solving benchmark CVRPs is compared. The population size of GA, PSM and VTPSO was 100; whereas, number of ants in ACO was equal to the number of nodes assigned to a vehicle as it desired. For the fair comparison, the number of iteration was set at 200 for the algorithms. The selected parameters are not optimal values, but considered for simplicity as well as for fairness in observation.

Table 5 compares CVRP cost for some standard Sweep based existing clustering methods and adaptive Sweep on A-VRP benchmark problems. Bottom of the table shows average and pairwise win/draw/loss summary over the total 27 problems. In adaptive Sweep, cluster formation starting angle is selected through proposed heuristic approach and it is -problem dependent. From the results presented in Table 5, it is watched that the greater part of the cases adaptive Sweep outperformed the exiting standard Sweep based clustering technique. The outperformance of adaptive Sweep is only for different starting angles in adaptive sweep for a particular route optimization, As an example, for n33-k6 problem, HHA achieved CVRP cost of 919. For the same problem the outcome of adaptive Sweep with adaptively selected starting angle 303.18° is 751. The results of route optimization with GA, ACO, PSM and VTPSO on adaptive Sweep cluster are compared with other existing methods. In most of the

cases adaptive sweep based methods outperformed corresponding standard Sweep based methods. For example adaptive sweep + PSM wins in 27, 9, 4 and 16 out of 27 cases, respectively. Only a few cases, standard Sweep based methods are found better than adaptive Sweep. For example Centroid based 3-phase and Sweep + Cluster Adjust outperforms Adaptive sweep for some problems. Adaptive Sweep is outperformed over other methods on the basis of average CVRP cost over 27 problems. The average CVRP cost for standard Sweep based methods are 1310.11, 1134.67, and 1181.44, respectively. On the other hand, achieved average CVRP cost for adaptive Sweep with GA, ACO, PSM and VTPSO are 1169.48, 1195.33, 1169.19 and 1168.63, respectively. Among adaptive Sweep based methods, PSM and VTPSO outperformed GA and ACO. Finally, CVRP cost with VTPSO are found best among the methods and it showed best (i.e., minimum) CVRP costs for all 27 problems.

Table 6 compares CVRP cost for some standard Sweep based existing clustering methods and adaptive Sweep on P-VRP benchmark problems. Bottom of the table shows average and pairwise win/draw/loss summary over the total 24 problems. From the table it is remarkable that most of the cases adaptive Sweep outperformed its corresponding other existing methods. A notable difference in the P-VRP problems from A-VRP problems of Table 5 is that adaptively selected starting angle is same for several problems. Examination on the data revealed that such problems have similar type of node coordinates. Adaptive Sweep is outperformed over existing optimization method on the basis of average

CVRP cost over 24 problems. The average CVRP cost for the methods are 708.33, 639.00, and 640.33 respectively. On the other hand, achieved average CVRP cost for adaptive Sweep with GA, ACO, PSM and VTPSO are 645.54, 655.71, 643.42 and

637.17, respectively. The routes obtained with GA, ACO, PSM and VTPSO on adaptive Sweep cluster outperformed corresponding other methods in 23, 12, 13, 6, 24 and 3 out of 24 cases, respectively.

Table 5: Comparison of CVRP cost for clustering with existing methods and Adaptive Sweep + SI methods on A-VRP benchmark problems

SI.	Problem	CVRP Cost for Existing methods			Starting Angle (Θs)	CVRP Cost for Adaptive Sweep + SI			
		HHA ¹⁰	Centroid-based 3-phase ⁸	Sweep + Cluster Adjustment ⁸		GA	ACO	PSM	VTPSO
1	n32-k5	1012	881	872	152.02	882	897	882	882
2	n33-k5	847	728	788	195.95	698	717	698	698
3	n33-k6	919	770	829	303.18	751	758	751	751
4	n34-k5	933	812	852	203.2	785	808	785	785
5	n36-k5	1126	814	884	323.13	881	917	881	881
6	n37-k5	876	756	734	248.84	756	774	756	754
7	n37-k6	1180	1027	1050	264.29	1112	1128	1112	1112
8	n38-k5	920	819	874	148.57	813	845	813	813
9	n39-k5	1147	864	971	180	877	914	877	877
10	n39-k6	1065	881	966	246.8	978	975	972	972
11	n44-k6	1356	1037	1092	253.3	1057	1116	1056	1056
12	n45-k6	1210	1040	1043	138.01	1075	1081	1073	1073
13	n45-k7	1361	1288	1281	180	1307	1339	1305	1305
14	n46-k7	1071	992	1013	75.96	977	1010	975	975
15	n48-k7	1292	1145	1143	3.18	1153	1165	1152	1152
16	n53-k7	1261	1117	1116	220.6	1091	1131	1090	1090
17	n54-k7	1414	1209	1320	4.09	1380	1374	1361	1361
18	n55-k9	1317	1155	1192	318.96	1191	1192	1190	1190
19	n60-k9	1733	1430	1574	170.54	1503	1528	1505	1503
20	n61-k9	1285	1201	1184	333.43	1170	1186	1164	1164
21	n62-k8	1604	1470	1559	263.66	1409	1435	1409	1408
22	n63-k9	2001	1766	1823	153.43	1824	1852	1823	1823
23	n63-k10	1542	1405	1523	6.34	1477	1511	1480	1477
24	n64-k9	1821	1587	1597	94.57	1598	1628	1598	1598
25	n65-k9	1429	1276	1351	237.99	1320	1327	1317	1317
26	n69-k9	1333	1283	1254	352.09	1269	1275	1259	1259
27	n80-k10	2318	1883	2014	149.04	2137	2195	2136	2136
Average		1310.11	1134.67	1181.44		1165.59	1188.07	1163.7	1163.41

Pairwise Win/Draw/Lose Summary

HHA -		27/0/0	27/0/0	-	27/0/0	27/0/0	27/0/0	27/0/0
Centroid-based 3-phase		-	7/0/20	-	9/1/17	5/0/22	9/1/17	10/0/17
Sweep -Algorithm + Cluster Adjust.		-	- -		16/0/11	15/1/11	14/1/12	15/1/11
GA -		-	- -			2/0/25	16/11/0	17/10/0
ACO -		-	- -				27/0/0	27/0/0
PSM -		-	- -					4/24/0

Table 6: Comparison of CVRP cost for clustering with existing methods and Adaptive Sweep + SI on P-VRP benchmark problems

SI.	Problem	CVRP Cost for Existing methods			Starting Angle (Θs)	CVRP Cost for Adaptive Sweep + SI			
		HHA ¹⁰	Centroid-based 3-phase ⁸	Sweep + Cluster Adjust ⁸		GA	ACO	PSM	VTPSO
1	n16-k8	546	497	568	335.1	549	554	549	549
2	n19-k2	253	256	236	335.1	246	246	246	246
3	n20-k2	267	240	238	335.1	249	249	249	249
4	n21-k2	288	240	238	335.1	211	217	211	211
5	n22-k2	274	245	237	335.1	216	223	216	216
6	n22-k8	667	672	687	238.39	633	633	633	633
7	n23-k8	743	703	645	333.43	634	636	634	634
8	n40-k5	563	505	499	119.48	492	504	483	483
9	n45-k5	662	533	525	119.48	524	556	524	524
10	n50-k7	647	583	585	278.43	589	599	583	583
11	n50-k8	721	669	675	278.43	677	713	677	677
12	n50-k10	808	740	779	278.43	783	793	783	783
13	n51-k10	857	779	806	208.3	804	822	802	802
14	n55-k7	679	610	611	278.43	595	619	595	595
15	n55-k8	690	654	601	242.59	589	608	589	586
16	n55-k10	832	749	763	278.43	745	767	745	745
17	n55-k15	1180	1022	1056	278.43	1099	1106	1099	1099
18	n60-k10	896	786	823	278.43	830	848	830	830
19	n60-k15	1159	1006	1086	278.43	1119	1136	1119	1119
20	n65-k10	964	836	856	278.43	859	880	859	859
21	n70-k10	989	891	902	278.43	914	951	911	911
22	n76-k4	753	685	603	104.04	681	630	658	612
23	n76-k5	671	737	647	144.16	689	675	662	647
24	n101-k4	891	698	702	115.46	766	772	785	699
	Average	708.33	639.00	640.33		645.54	655.71	643.42	637.17

Pairwise Win/Draw/Lose Summary

Hybrid Heuristic	-	21/0/3	22/0/2	-	22/0/2	22/0/2	22/0/2	23/0/1
Centroid-based 3-phase	-	-	-10/0/14	-	13/0/11	9/0/15	12/1/11	12/1/11
Sweep + Cluster Adjust.	-	-	-	-	-11/0/13	6/0/18	12/0/12	13/1/10
GA	-	-	-	-	-	2/3/19	6/17/1	6/16/0
ACO	-	-	-	-	-	-	22/0/2	24/0/0
PSM	-	-	-	-	-	-	-	3/21/0

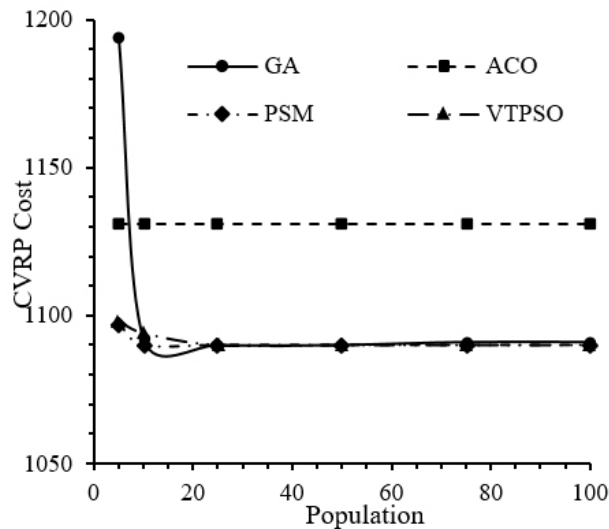
Experimental Analysis

This section investigates the effect of population size in route optimizing phase as the selected methods work with a populations of solutions. For GA, PSM and VTPSO, population size was varied from 5 to 100 whereas, in ACO the number of ants equals the

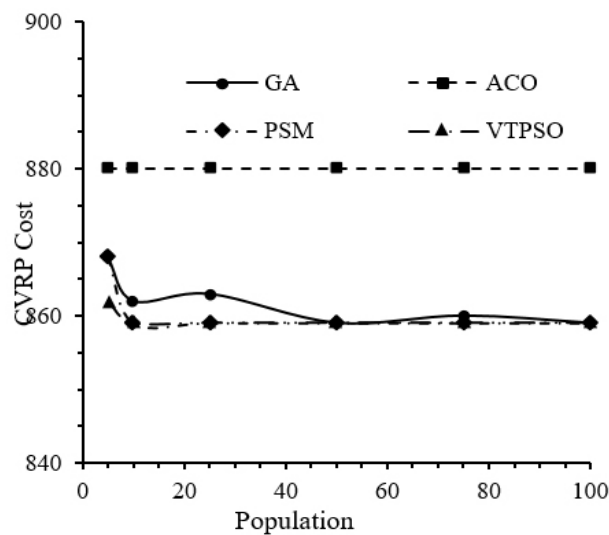
number of cities as it desired. Fig. 3 shows CVRP cost (i.e., total travel cost) for population variation; experiment conducted for fixed 200 iteration for fair comparison. The number of clusters (i.e., vehicles) were 7 and 10 for A: n53-k7 and P: n65-k10 problems, respectively. From the figure it is observed

that CVRP cost is invariant for ACO because population variation was not employed for it. On the other hand, GA is most sensitive with population size: CVRP cost through GA was very bad with respect to others at small population size (e.g., 5) and was competitive at larger population size. From the figure

it is also observed that recent SI methods PSM and VTPSO are better than ACO and GA in population variation. At a glance, VTPSO is shown to outperform any other method for any population size and PSM is competitive to VTPSO.



(a) A: n53-k7 problem



(b) P :n65-k10 problem

Fig. 3: CVRP cost (i.e., total travel cost) for population variation

Conclusions

A popular way of solving CVRP is to cluster the nodes according to vehicles using Sweep algorithm first and then generate route for each vehicle with TSP algorithm. In this study adaptive Sweep method is considered for appropriate cluster formation. Finally, GA, ACO, PSM and VTPSO are applied

to generate optimal route with the clusters. The experimental results on the benchmark problems revealed that heuristically selected starting angles have positive effect on Sweep clustering and adaptive Sweep plus SI methods is an effective CVRP solving method when compared with several related existing methods.

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